

Characterizing Bit Permutation Networks

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Received August 1997; accepted October 1997

Abstract: In recent years, many multistage interconnection networks using 2×2 switching elements have been proposed for parallel architectures. Typical examples are baseline networks, banyan networks, shuffle-exchange networks, and their inverses. As these networks are blocking, such networks with extra stages have also been studied extensively. These include Benes networks and $\Delta \oplus \Delta'$ networks. Recently, Hwang et al. studied k -extra-stage networks, which are a generalization of the above networks. They also investigated the equivalence issue among some of these networks. In this paper, we studied a more general class of networks, which we call $(m + 1)$ -stage d -nary bit permutation networks. We characterize the equivalence of such networks by sequence of positive integers. © 1999 John Wiley & Sons, Inc. Networks 33: 261–267, 1999

Keywords: multistage interconnection network; switching network; permutation routing; Sterling number; rearrangeably nonblocking

1. INTRODUCTION

Consider a multistage interconnection network $\mathcal{V}$ with $N = d^{n+1}$ inputs and outputs and which has $m + 1$ stages of N/d crossbars of size $d \times d$. Let the j th crossbar in a stage be labeled by j in the d -nary number (with n bits). A *bit- i group* consists of those crossbars whose labels are identical except the i th bit. Such a group will be labeled by a d -nary number x with n bits which is identical to any member in the group except that bit i is replaced by the symbol x_0 , which stands for the set $\{0, 1, \dots, d - 1\}$. $\mathcal{V}$ will be called a $(m + 1)$ -stage d -nary bit permutation network if the linking between stage k to stage $k + 1$ is always from a bit- i_k group G to a bit- j_k group G' , where G' is a permutation of G , for $k = 0, 1, \dots, m - 1$.

For a detailed description and notation of bit permutation networks, see Section 2.

Note that $(n + 1)$ -stage binary bit permutation networks include all self-routing networks like Omega, banyan, baseline, and their inverse networks. Binary bit permutation networks have been widely studied in the literature [1, 3, 6, 8, 11] for their topological equivalence. Bermond et al. [3] characterized the Omega-equivalent class by the $P(i, j)$ property. An $(n + 1)$ -stage network satisfies the $P(i, j)$ property if the subnetwork from stage i to stage j has exactly 2^{n-j+i} components. Then, an $(n + 1)$ -stage network with the unique path property is in the Omega-equivalent class if and only if it satisfies the $P(i, j)$ property for all $0 \leq i \leq j \leq n$.

Another special class of bit permutation networks consists of $(n + 1)$ -stage networks with extra stages. A k -extra-stage network is a cascade of a $(n + 1)$ -stage network with k extra stages also satisfying the bit permutation linking pattern. Lea and Shyy [7, 9] proposed adding extra stages to a binary inverse banyan network while the

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Contract grant sponsor: National Science Council; contract grant number: NSC86-2115-M0009-002

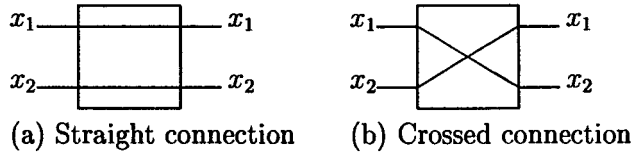


Fig. 1. A 2×2 switching element.

k extra stages are added by pattern F^{-1} (see below). Hwang et al. [5] generalized the study of equivalence by adding extra stages to a binary Omega-equivalent network with the following patterns for extra stages:

- (i) F : They are identical to the first k stages of the network;
- (ii) F^{-1} : Identical to the mirror image of the first k stages;
- (iii) L : Identical to the last k stages;
- (iv) L^{-1} : Identical to the mirror image of the last k stages.

In this paper, we determined the equivalence classes among all $(m + 1)$ -stage d -nary bit permutation networks. We characterize such a network by an m -sequence over $\{1, 2, \dots, n\}$, namely, every $(m + 1)$ -stage d -nary bit permutation network is reduced to an m -sequence over $\{1, 2, \dots, n\}$ and equivalence is determined by some easily computable sequence statistics. Note that the sequence is independent of d . For $m = n$, this characterization, of course, corresponds to the $P(i, j)$ characterization. But the sequence-graph correspondence is not in an obvious way. With the power and convenience of the sequence characterization, we easily give an explicit solution of the size of the s -stage bit permutation class. Recently, Hu et al. [4] gave an $O(N^4 \log N)$ -time algorithm to check the equivalence of combined $(2n - 1)$ -stage networks, which are obtained by cascading two Omega-equivalent networks. We give an (mn) -time algorithm for checking the equivalence of two $(m + 1)$ -stage bit permutation networks. In particular, the running time is $O(\log^2 N)$ when the network has $2n - 1$ stages.

2. NETWORKS

We start the discussion of bit permutation networks by examining the following classical example: A typical Omega-equivalent network consists of N input terminals, N output terminals, and $\log_2 N$ columns (stages) of 2×2 switching elements in which each column has $N/2$ switching elements. Figure 1 shows a 2×2 switching

element, and Figure 2 shows a baseline network with $N = 16$, in which a terminal i is represented by its binary number representation (x_3, x_2, x_1, x_0) and is adjacent to a switching element named by (x_3, x_2, x_1) .

One can view the baseline network in Figure 2 as a graph whose vertices are those 32 switching elements named by $(x_3, x_2, x_1)_i$, where $0 \leq i \leq 3$ and $x_1, x_2, x_3 \in \{0, 1\}$, and there are links

- from $(x_3, x_2, x_1)_0$ to $(x_0, x_3, x_2)_1$,
- from $(x_3, x_2, x_1)_1$ to $(x_3, x_0, x_2)_2$, and
- from $(x_3, x_2, x_1)_2$ to $(x_3, x_2, x_0)_3$,

where $x_0 \in \{0, 1\}$, meaning there are two links, one with $x_0 = 0$ and the other with $x_0 = 1$. A link from a switching element x at stage i to a switching element y at stage $i + 1$ exists if the bits of y can be obtained from the bits of x by a permutation depending only on i . For instance, we can represent the links from stage 0 to stage 1 by a permutation f_1 on $\{0, 1, 2, 3\}$ with

$$f_1(0) = 1, \quad f_1(1) = 2, \quad f_1(2) = 3, \quad f_1(3) = 0.$$

In this way, the links from stage 0 to stage 1 are those from $(x_3, x_2, x_1)_0$ to $(x_{f_1(3)}, x_{f_1(2)}, x_{f_1(1)})_1$. We can also say that for a link from $x = (x_3, x_2, x_1)_0$ to $y = (x_0, x_3, x_1)_1$ a coordinate x_j at the j th position of x moves to the $f_1^{-1}(j)$ th position of y , where the coordinate x_0 's moving from "outside" into y means there are two such links. Similarly, the following permutations f_2 and f_3 represent links from stage 1 to stage 2 and stage 2 to stage 3, respectively:

$$\begin{aligned} f_2(0) &= 1, & f_2(1) &= 2, & f_2(2) &= 0, & f_2(3) &= 3; \\ f_3(0) &= 1, & f_3(1) &= 0, & f_3(2) &= 2, & f_3(3) &= 3. \end{aligned}$$

Throughout this paper, we shall use the *cycle* notation for

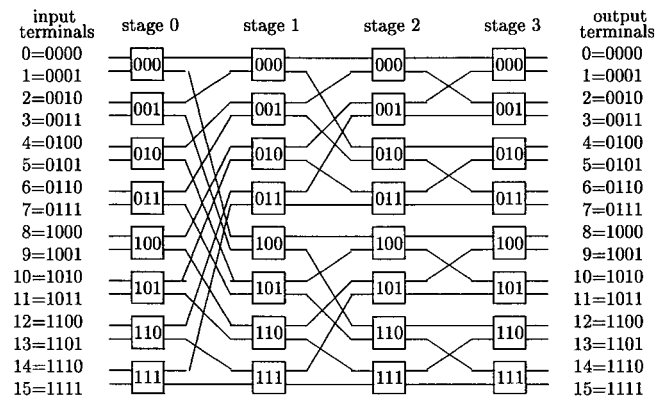


Fig. 2. A baseline network with $N = 16$.

permutations, that is, the cycle $(i_1, i_2, \dots, i_n)$ represents the permutation f with

$$f(i_1) = i_2, \quad f(i_2) = i_3, \dots, \quad f(i_{n-1}) = i_n, \\ f(i_n) = i_1, \quad \text{and } f(j) = j \text{ for all other } j.$$

Then, f_1 can be represented by $(0, 1, 2, 3)$, f_2 by $(0, 1, 2)$, and f_3 by $(0, 1)$.

Not only can the baseline network in Figure 2 be represented by the permutations f_i , but also all bit permutation networks can be represented in this way. We can also use general $d \times d$ switching elements instead of 2×2 switching elements. A general setting is as follows:

Definition. Suppose that n is a positive integer and $f_1, f_2, \dots, f_m$ are $m \geq 0$ permutations on $\{0, 1, \dots, n\}$ such that $f_i(0) \neq 0$ for $1 \leq i \leq m$. The $(m+1)$ -stage d -nary bit permutation network $N_d(n, f_1, f_2, \dots, f_m)$ is the network whose vertices are those $(x_n, x_{n-1}, \dots, x_1)_i$ with $0 \leq i \leq m$ and $x_j \in \{0, 1, \dots, d-1\}$ for $1 \leq j \leq n$, and each $(x_n, x_{n-1}, \dots, x_1)_{i-1}$ is adjacent to $(x_{f_i(n)}, x_{f_i(n-1)}, \dots, x_{f_i(1)})_i$, where $x_0 \in \{0, 1, \dots, d-1\}$. In other words, there is an edge from $(x_n, x_{n-1}, \dots, x_1)_{i-1}$ to $(y_n, y_{n-1}, \dots, y_1)_i$ whenever $y_j = x_{f_i(j)}$ for $1 \leq j \leq n$.

The following examples show $(m+1)$ -stage binary bit permutation networks with $N = 2^{n+1}$ input and output terminals. Examples 1–3 have been worked out by Wu and Feng [11]. We give the presentation in our format for easier use later.

Example 1. The baseline network BL is precisely $N_2(n, f_1, f_2, \dots, f_n)$, with

$$f_i = (0, 1, \dots, n-i+1) \quad \text{for } 1 \leq i \leq n.$$

The inverse baseline network BL^{-1} is $N_2(n, f_1, f_2, \dots, f_n)$, with

$$f_i = f_{n-i+1}^{-1} = (i, i-1, \dots, 1, 0) \quad \text{for } 1 \leq i \leq n.$$

Example 2. The banyan network BY or the indirect binary n -cube network is $N_2(n, f_1, f_2, \dots, f_n)$, with

$$f_i = (0, i) \quad \text{for } 1 \leq i \leq n.$$

The modified data manipulator is $N_2(n, f_1, f_2, \dots, f_n)$, with

$$f_i = f_{n-i+1}^{-1} = (0, n-i+1) \quad \text{for } 1 \leq i \leq n$$

and, hence, is the inverse banyan network BY^{-1} .

Example 3. The Omega (or shuffle-exchange) network SE is $N_2(n, f_1, f_2, \dots, f_n)$, with

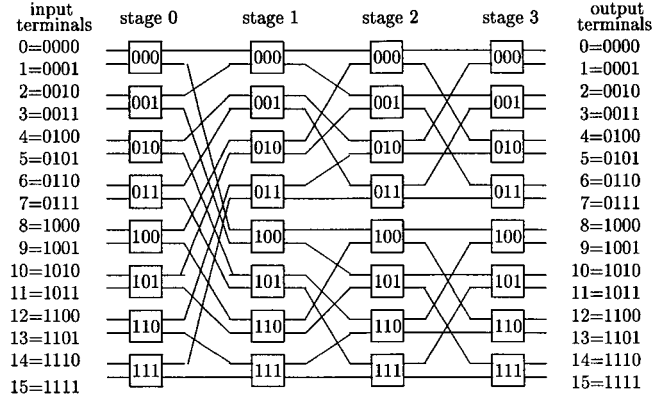


Fig. 3. The network in Example 4.

$$f_i = (n, n-1, \dots, 1, 0) \quad \text{for } 1 \leq i \leq n.$$

The flip (or inverse-shuffle-exchange) network SE^{-1} is $N_2(n, f_1, f_2, \dots, f_n)$, with

$$f_i = f_i^{-1} = (0, 1, \dots, n) \quad \text{for } 1 \leq i \leq n.$$

Example 4. Interchange the role of the coordinates x_1 and x_2 of (x_3, x_2, x_1) in stages 2 and 3 of the baseline network in Figure 2. The resulting network, which is topologically equivalent to BL , is $N_2(3, f_1, f_2, f_3)$ with

$$f_1 = (0, 1, 2, 3), \quad f_2 = (0, 2, 1), \quad \text{and } f_3 = (0, 2)$$

(see Fig. 3). Note that if one interchanges the switching elements $(x_3, 0, 1)_i$ with $(x_3, 1, 0)_i$ for $i = 2, 3$ the resulting network is the baseline network.

Example 5. Suppose that $N_2(n, f_1, f_2, \dots, f_n)$ is an Omega-equivalent network Δ and $N_2(n, f'_1, f'_2, \dots, f'_n)$ is another Omega-equivalent network Δ' . Then, $\Delta \oplus \Delta'$ is $N_2(n, f_1, f_2, \dots, f_n, f'_1, f'_2, \dots, f'_n)$.

Example 6. Suppose that $N_2(n, f_1, f_2, \dots, f_n)$ is an Omega-equivalent network Δ and $1 \leq k \leq n$. Hwang et al. [5] defined four k -extra-stage networks as follows: $\Delta_F(k)$ [respectively, $\Delta_L(k)$] is Δ together with k extra stages identical to the (respectively, last) first k stages of Δ , and $\Delta_F^{-1}(k)$ [respectively, $\Delta_L^{-1}(k)$] is Δ together with k extra stages identical to the mirror image of the first (respectively, last) k stages of Δ . Then,

$$\Delta_F(k) \text{ is } N_2(n, f_1, f_2, \dots, f_n, f_1, f_2, \dots, f_k);$$

$$\Delta_F^{-1}(k) \text{ is } N_2(n, f_1, f_2, \dots, f_n, f_k^{-1}, f_{k-1}^{-1}, \dots, f_1^{-1});$$

$$\Delta_L(k) \text{ is } N_2(n, f_1, f_2, \dots, f_n, f_{n-k+1}, f_{n-k+2}, \dots, f_n);$$

$\Delta_L^{-1}(k)$ is $N_2(n, f_1, f_2, \dots, f_n, f_n^{-1}, f_{n-1}^{-1}, \dots, f_{n-k+1}^{-1})$.

3. TOPOLOGICAL EQUIVALENCE

The main effort of this paper was to establish methods for determining if two bit permutation networks $N_d(n, f_1, f_2, \dots, f_m)$ and $N_d(n, f'_1, f'_2, \dots, f'_m)$ are topologically equivalent or graphically isomorphic. The following theorem is the foundation of our theory:

Theorem 1. *If there exist permutations $g_0, g_1, \dots, g_m$ on $\{0, 1, \dots, n\}$ such that $g_i(0) = 0$ for $0 \leq i \leq m$ and $f'_i = g_{i-1}^{-1} \circ f_i \circ g_i$ for $1 \leq i \leq m$, then $N_d(n, f_1, f_2, \dots, f_m)$ is isomorphic to $N_d(n, f'_1, f'_2, \dots, f'_m)$.*

Proof. Consider the bijection g from the vertex set of $N_d(n, f_1, f_2, \dots, f_m)$ to the vertex set of $N_d(n, f'_1, f'_2, \dots, f'_m)$ defined by

$$(x_n, x_{n-1}, \dots, x_1)_i \xrightarrow{g} (x_{g_i(n)}, x_{g_i(n-1)}, \dots, x_{g_i(1)})_i$$

for $0 \leq i \leq m$.

In other words, $g((x_n, x_{n-1}, \dots, x_1)_i) = (x'_n, x'_{n-1}, \dots, x'_1)_i$ whenever $x'_j = x_{g_i(j)}$ for $1 \leq j \leq n$.

To see that these two networks are isomorphic, we only need to check that g is edge-preserving. Suppose that e is an edge from $(x_n, x_{n-1}, \dots, x_1)_{i-1}$ to $(y_n, y_{n-1}, \dots, y_1)_i$ in $N_d(n, f_1, f_2, \dots, f_m)$, that is, $y_j = x_{f_i(j)}$ for $1 \leq j \leq n$. Let

$$g((x_n, x_{n-1}, \dots, x_1)_{i-1}) = (x'_n, x'_{n-1}, \dots, x'_1)_{i-1},$$

i.e., $x'_j = x_{g_{i-1}(j)}$ for $1 \leq j \leq n$,

and

$$g((y_n, y_{n-1}, \dots, y_1)_i) = (y'_n, y'_{n-1}, \dots, y'_1)_i,$$

i.e., $y'_j = y_{g_i(j)}$ for $1 \leq j \leq n$.

Then,

$$\begin{aligned} y'_j &= y_{g_i(j)} = x_{f_i \circ g_i(j)} = x_{g_{i-1} \circ f'_i(j)} \\ &= x'_{f'_i(j)} \quad \text{for } 1 \leq j \leq n. \end{aligned}$$

Thus, there exists an edge from $(x'_n, x'_{n-1}, \dots, x'_1)_{i-1}$ to $(y'_n, y'_{n-1}, \dots, y'_1)_i$. Conversely, an edge in $N_d(n, f'_1, f'_2, \dots, f'_m)$ also corresponds to an edge in $N_d(n, f_1, f_2, \dots, f_m)$. ■

As a quick application of Theorem 1, consider BL

in Figure 2 as $N_2(3, f_1, f_2, f_3)$, with $f_1 = (0, 1, 2, 3)$, $f_2 = (0, 1, 2)$, and $f_3 = (0, 1)$ and the network in Example 4 as $N_2(3, f'_1, f'_2, f'_3)$, with $f'_1 = (0, 1, 2, 3)$, $f'_2 = (0, 2, 1)$, and $f'_3 = (0, 2)$. Then, $g_0 = (2, 3)$ and $g_1 = g_2 = g_3 = (1, 2)$ show that these two networks are isomorphic.

Theorem 2. *Every bit permutation network $N_d(n, f_1, f_2, \dots, f_m)$ is isomorphic to some bit permutation network $N_d(n, f'_1, f'_2, \dots, f'_m)$, where $f'_i = (0, k_i)$ with $k_i \in \{1, 2, \dots, n\}$ for $1 \leq i \leq m$.*

Proof. We shall prove by induction on j the following claim which implies the theorem:

Claim(j). $N_d(n, f_1, f_2, \dots, f_m)$ is isomorphic to some $N_d(n, f'_1, f'_2, \dots, f'_m)$, where $f'_i = (0, k_i)$ with $k_i \in \{1, 2, \dots, n\}$ for $1 \leq i \leq j$.

Claim(0) is clearly true. Suppose that Claim($j - 1$) holds. Consider a general $j \geq 1$. Let $k_j = f'_j(0)$. Then, $k_j \in \{1, 2, \dots, n\}$. Let $g_j = (f'_j)^{-1} \circ (0, k_j)$ and all other g_i are identity permutations. It is easy to check that $g_i(0) = 0$ for $0 \leq i \leq m$. According to Theorem 1, $N_d(n, f'_1, f'_2, \dots, f'_m)$, and, therefore, $N_d(n, f_1, f_2, \dots, f_m)$, is isomorphic to $N_d(n, f''_1, f''_2, \dots, f''_m)$, where $f''_i = g_{i-1}^{-1} \circ f'_i \circ g_i$ for $1 \leq i \leq m$. In particular, $f''_i = f'_i = (0, k_i)$ for $1 \leq i \leq j - 1$, and $f''_j = f'_j \circ g_j = f'_j \circ (f'_j)^{-1} \circ (0, k_j) = (0, k_j)$. This gives Claim(j). ■

So, basically, we only need to consider the networks with $f_i = (0, k_i)$. For convenience, we shall use $N_d(n, k_1, k_2, \dots, k_m)$ as a short notation for the network $N_d(n, f_1, f_2, \dots, f_m)$, with $f_i = (0, k_i)$ and $k_i \in \{1, 2, \dots, n\}$ for $1 \leq i \leq m$.

Our next step is to determine when two networks are isomorphic.

Theorem 3. *If g is a permutation on $\{0, 1, 2, \dots, m\}$ with $g(0) = 0$, then the network $N_d(n, k_1, k_2, \dots, k_m)$ is isomorphic to $N_d(n, g(k_1), g(k_2), \dots, g(k_m))$.*

Proof. The theorem following from Theorem 1 and the fact that $g \circ (0, k_i) \circ g^{-1} = (0, g(k_i))$. Note that we apply Theorem 1 by choosing all g_i as g . ■

For any sequence $(k_1, k_2, \dots, k_m)$ over $\{1, 2, \dots, n\}$ with $a = |\{k_1, k_2, \dots, k_m\}|$, let $1 = i_1 < i_2 < \dots < i_a$ be those indices such that $k_{i_r} \notin \{k_1, k_2, \dots, k_{i_r-1}\}$ for $1 \leq r \leq a$. Choose any permutation g on $\{1, 2, \dots, n\}$ with the property that $g(k_{i_r}) = r$ for $1 \leq r \leq a$. Then, $(k_1^*, k_2^*, \dots, k_m^*) = (g(k_1), g(k_2), \dots, g(k_m))$ is a sequence with the property that $\{k_1^*, k_2^*, \dots, k_m^*\} = \{1, 2, \dots, a\}$ and for any $1 \leq r \leq a$, $\{k_1^*, k_2^*, \dots, k_{i_r}^*\} = \{1, 2, \dots, r\}$ when i_r is the mini-

mum index with $k_r^* = r$. Such a sequence is called a *canonical sequence* over $\{1, 2, \dots, n\}$.

Corollary 4. Any $N_d(n, f_1, f_2, \dots, f_m)$ is isomorphic to $N_d(n, k_1, k_2, \dots, k_m)$ for some canonical sequence over $\{1, 2, \dots, n\}$.

Theorem 5. If $a = |\{k_1, k_2, \dots, k_m\}|$, then $N_d(n, k_1, k_2, \dots, k_m)$ has d^{n-a} connected components.

Proof. First, we claim that for any two vertices $x = (x_n, x_{n-1}, \dots, x_1)_i$ and $y = (y_n, y_{n-1}, y_{n-1}, \dots, y_1)_j$ such that $x_k = y_k$ for all $k \in \{k_1, k_2, \dots, k_m\}$ there exists a path joining them.

For the case in which $i = 0$ and $j = m$, $x = x^{(0)} \rightarrow x^{(1)} \rightarrow \dots \rightarrow x^{(m)} = y$ is a desired path, where $x^{(i)}$ is a vertex in stage i that is obtained from $x^{(i-1)}$ by replacing its k_i -th coordinate with y_{k_i} . For the general case, consider the vertices $x^* = (x_n, x_{n-1}, \dots, x_1)_0$ and $y^* = (y_n, y_{n-1}, \dots, y_1)_m$. By the above construction, there exists an $x^* - y^*$ path. It is clear that $(x_n, x_{n-1}, \dots, x_1)_i \rightarrow (x_n, x_{n-1}, \dots, x_1)_{i-1} \rightarrow \dots \rightarrow (x_n, x_{n-1}, \dots, x_1)_0 = x^*$ is an $x - x^*$ path. Similarly, there exists a $y^* - y$ path. Thus, x and y are joined by a path.

On the other hand, a move from any vertex to its neighbor never changes its k -th coordinate for $k \notin \{k_1, k_2, \dots, k_m\}$. Thus, two vertices with different values at the k -th coordinate for some $k \notin \{k_1, k_2, \dots, k_m\}$ are not in the same connected component of the network. Therefore, the network has d^{n-a} connected components. ■

Theorem 6. For any two canonical sequences $(k_1, k_2, \dots, k_m)$ and $(k'_1, k'_2, \dots, k'_m)$ over $\{1, 2, \dots, n\}$, the network $N_d(n, k_1, k_2, \dots, k_m)$ is isomorphic to $N_d(n, k'_1, k'_2, \dots, k'_m)$ if and only if $(k_1, k_2, \dots, k_m) = (k'_1, k'_2, \dots, k'_m)$.

Proof. The two networks are clearly isomorphic if the two sequences are equal.

Conversely, suppose that $(k_1, k_2, \dots, k_m) \neq (k'_1, k'_2, \dots, k'_m)$. Let i be the minimum index such that $k_i \neq k'_i$, say $k_i < k'_i$ and $k_r = k'_r$ for $1 \leq r \leq i - 1$. Since the two sequences are canonical, $i \geq 2$. By the definition, there exists some $1 \leq r \leq i - 1$ such that $k_r = k'_r \in \{k_i, k'_i\}$. Choose a maximum such r . Then, $|\{k_r, k_{r+1}, \dots, k_i\}| = a \neq a' = |\{k'_r, k'_{r+1}, \dots, k'_i\}|$. Note that $N_d(n, k_r, k_{r+1}, \dots, k_i)$ has d^{m-a} connected components while $N_d(n, k'_r, k'_{r+1}, \dots, k'_i)$ has $d^{m-a'}$ connected components. Hence, the two networks are not isomorphic. ■

Remark. By Theorem 6, the $N_d(n, k_1, k_2, \dots, k_m)$ in Corollary 4 is unique. Such a network is called the *canonical representation* of $N_d(n, f_1, f_2, \dots, f_m)$.

A condition slightly weaker than the $P(i, j)$ characterization follows from Theorem 6:

Corollary 7. An $(n + 1)$ -stage network with the unique path property and satisfying the $P(i, i + 1)$ property for $i = 0, 1, \dots, n$, is in the Omega-equivalent class if and only if it also satisfies the $P(0, j)$ property for $j = 1, 2, \dots, n$.

Proof. By satisfying the $P(i, i + 1)$ property, the network is a bit permutation network. By satisfying the $P(0, j)$ property, the number of components is increasing in j , which implies that $k_j \neq k_i$ for all $i < j$. It follows that $(k_1, k_2, \dots, k_n)$ are all distinct, but there is only one such canonical sequence. Corollary 7 follows from Theorem 6 immediately. ■

Theorem 8. The number of equivalent classes among $(m + 1)$ -stage d -nary bit permutation networks is

$$\sum_{t=1}^n \sum_{i=0}^t \frac{1}{t!} (-1)^{t-i} \binom{t}{i} i^m.$$

Proof. By Corollary 4 and Theorem 6, we only need to count the set $C(m, t)$ of canonical sequences of length m with t distinct elements for $1 \leq t \leq n$. Denote by $P(m, t)$ the set of all partitions of $\{1, 2, \dots, m\}$ into t nonempty subsets. Define a mapping h from $C(m, t)$ to $P(m, t)$ by $h(k_1, k_2, \dots, k_m) = \{C_1, C_2, \dots, C_t\}$, where $C_i = \{j : k_j = i\}$. It is easy to see that h is one-to-one. On the other hand, for any partition $\{C_1, C_2, \dots, C_t\}$ in $P(m, t)$, we may assume that $\min C_1 \leq \min C_2 \leq \dots \leq \min C_t$. Let $k = (k_1, k_2, \dots, k_m)$, where $k_i = j$ for $i \in C_j$. Then, $h(k) = \{C_1, C_2, \dots, C_t\}$. So, h is onto. Therefore, $|C(m, t)| = |P(m, t)|$, which by definition is the Stirling number of the second kind $S(m, t)$ (see [10]). It is well known that

$$S(m, t) = \frac{1}{t!} \sum_{i=0}^t (-1)^{t-i} \binom{t}{i} i^m.$$

Summing over t , we obtain Theorem 8. ■

4. ALGORITHM AND APPLICATIONS

By the theorems in Section 3, a bit permutation network $N_d(n, f_1, f_2, \dots, f_m)$ is topologically equivalent to its canonical representation $N_d(n, k_1, k_2, \dots, k_m)$, whose topology is determined by the canonical sequence $(k_1, k_2, \dots, k_m)$. We shall summarize an efficient algorithm from the proofs of Theorems 2 and 3. We then apply it to determine the equivalence among the networks mentioned in the examples of Section 2.

Algorithm. Find the canonical representation of a bit permutation network.

Input. A bit permutation network $N_d(n, f_1, f_2, \dots, f_m)$.

Output. The canonical sequence $(k_1^*, k_2^*, \dots, k_m^*)$ for $N_d(n, f_1, f_2, \dots, f_m)$.

Method.

```

 $f_0 = (0)$ ;  $f_{m+1} = (0)$ ;
for  $j = 1$  to  $m$  do  $f_j = f_j$ ;
for  $j = 1$  to  $m$  do
   $k_j = f_j(0)$ ;
   $f_{j+1} = (0, k_j) \circ f_j \circ f_{j+1}$ ;
end for;
for  $j = 1$  to  $n$  do  $\text{mark}[j] = 0$ ;
 $a = 0$ ;
for  $j = 1$  to  $m$  do
  if  $\text{mark}[k_j] = 0$  then
     $a = a + 1$ ;
     $\text{mark}[k_j] = a$ ;
     $k_j^* = \text{mark}[k_j]$ ;
  end then;
end for;

```

The time complexity of this algorithm is $O(mn)$.

Now, we use the above algorithm to get the canonical sequences of k -extra-stage Omega-equivalent networks.

Example 7. The following are the canonical sequences of $\Delta_F(k)$ for the six different Omega-equivalent networks Δ :

- (a) For $\Delta = SE$, the sequence is $(1, 2, \dots, m, 1, 2, \dots, k)$.
- (b) For $\Delta = SE^{-1}$, the sequence is $(1, 2, \dots, m, 1, 2, \dots, k)$.
- (c) For $\Delta = BL$, the sequence is $(1, 2, \dots, m, m, m - 1, \dots, m - k + 1)$.
- (d) For $\Delta = BL^{-1}$, the sequence is $(1, 2, \dots, m, m, m - 1, \dots, m - k + 1)$.
- (e) For $\Delta = BY$, the sequence is $(1, 2, \dots, m, 1, 2, \dots, k)$.
- (f) For $\Delta = BY^{-1}$, the sequence is $(1, 2, \dots, m, 1, 2, \dots, k)$.

In the type $\Delta_F(k)$, there are two nonequivalent classes.

Example 8. The following are the canonical sequences of $\Delta_F^{-1}(k)$ for the six different Omega-equivalent networks Δ .

- (a) For $\Delta = SE$, the sequence is $(1, 2, \dots, m, m, m - 1, \dots, m - k + 1)$.
- (b) For $\Delta = SE^{-1}$, the sequence is $(1, 2, \dots, m, m, m - 1, \dots, m - k + 1)$.

- (c) For $\Delta = BL$, the sequence is $(1, 2, \dots, m, k, k - 1, \dots, 1)$.
- (d) For $\Delta = BL^{-1}$, the sequence is $(1, 2, \dots, m, m, m - 1, \dots, m - k + 1)$.
- (e) For $\Delta = BY$, the sequence is $(1, 2, \dots, m, k, k - 1, \dots, 1)$.
- (f) For $\Delta = BY^{-1}$, the sequence is $(1, 2, \dots, m, k, k - 1, \dots, 1)$.

In the type $\Delta_F^{-1}(k)$, there are two nonequivalent classes.

Example 9. The following are the canonical sequences of $\Delta_L(k)$ for the six different Omega-equivalent networks Δ .

- (a) For $\Delta = SE$, the sequence is $(1, 2, \dots, m, 1, 2, \dots, k)$.
- (b) For $\Delta = SE^{-1}$, the sequence is $(1, 2, \dots, m, 1, 2, \dots, k)$.
- (c) For $\Delta = BL$, the sequence is $(1, 2, \dots, m, m, m - 1, \dots, m - k + 1)$.
- (d) For $\Delta = BL^{-1}$, the sequence is $(1, 2, \dots, m, k, k - 1, \dots, 1)$.
- (e) For $\Delta = BY$, the sequence is $(1, 2, \dots, m, m - k + 1, m - k + 2, \dots, m)$.
- (f) For $\Delta = BY^{-1}$, the sequence is $(1, 2, \dots, m, m - k + 1, m - k + 2, \dots, m)$.

In the type $\Delta_L(k)$, there are four nonequivalent classes.

Example 10. The following are the canonical sequences of $\Delta_L^{-1}(k)$ for the six different Omega-equivalent networks Δ :

- (a) For $\Delta = SE$, the sequence is $(1, 2, \dots, m, m, m - 1, \dots, m - k + 1)$.
- (b) For $\Delta = SE^{-1}$, the sequence is $(1, 2, \dots, m, m, m - 1, \dots, m - k + 1)$.
- (c) For $\Delta = BL$, the sequence is $(1, 2, \dots, m, m, m - 1, \dots, m - k + 1)$.
- (d) For $\Delta = BL^{-1}$, the canonical sequence is $(1, 2, \dots, m, m, m - 1, \dots, m - k + 1)$.
- (e) For $\Delta = BY$, the sequence is $(1, 2, \dots, m, m, m - 1, \dots, m - k + 1)$.
- (f) For $\Delta = BY^{-1}$, the sequence is $(1, 2, \dots, m, m, m - 1, \dots, m - k + 1)$.

In the type $\Delta_L^{-1}(k)$, there is one equivalent class.

If $k = m$, then the k -extra-stage networks have only two nonequivalent classes among $\Delta_F(k)$, $\Delta_F^{-1}(k)$, $\Delta_L(k)$, and $\Delta_L^{-1}(k)$, namely, one characterized by se-

quence $(1, 2, \dots, m, 1, 2, \dots, m)$, and the other, by $(1, 2, \dots, m, m, m-1, \dots, 1)$. Since the Benes network $BL_L^{-1}(n)$ is in the second class, networks in the second class are rearrangeable. Benes [2] conjectured that $SE_F(n)$ is rearrangeable. If the conjecture is true, then the networks in the first class are also rearrangeable.

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