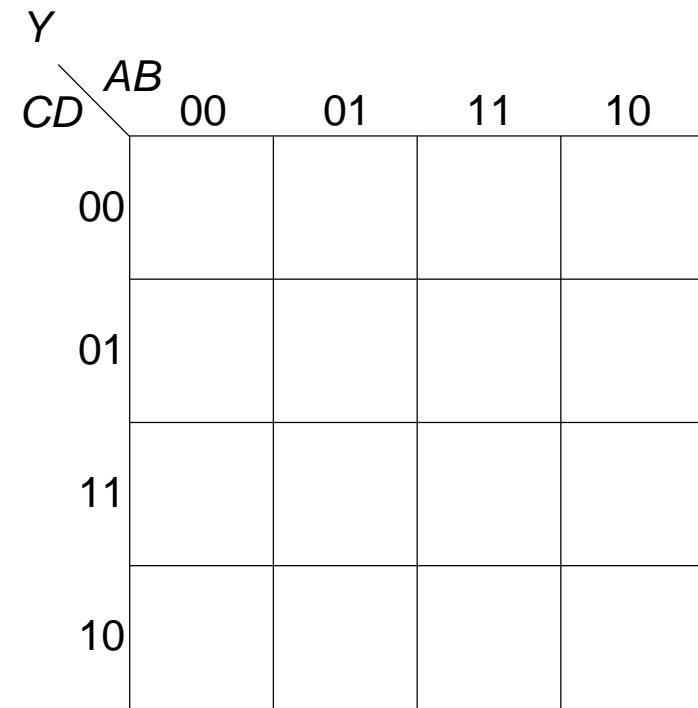


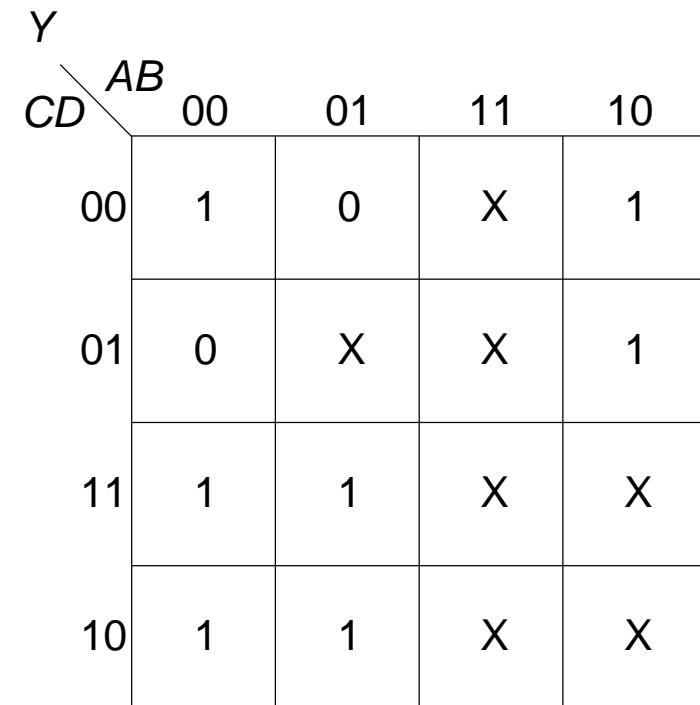
K-Maps with Don't Cares

A	B	C	D	Y
0	0	0	0	1
0	0	0	1	0
0	0	1	0	1
0	0	1	1	1
0	1	0	0	0
0	1	0	1	X
0	1	1	0	1
0	1	1	1	1
1	0	0	0	1
1	0	0	1	1
1	0	1	0	X
1	0	1	1	X
1	1	0	0	X
1	1	0	1	X
1	1	1	0	X
1	1	1	1	X



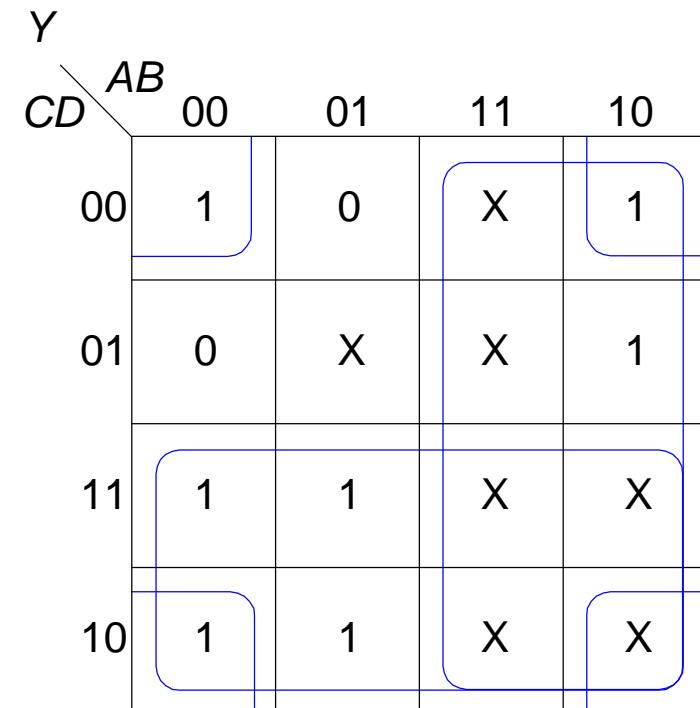
K-Maps with Don't Cares

A	B	C	D	Y
0	0	0	0	1
0	0	0	1	0
0	0	1	0	1
0	0	1	1	1
0	1	0	0	0
0	1	0	1	X
0	1	1	0	1
0	1	1	1	1
1	0	0	0	1
1	0	0	1	1
1	0	1	0	X
1	0	1	1	X
1	1	0	0	X
1	1	0	1	X
1	1	1	0	X
1	1	1	1	X



K-Maps with Don't Cares

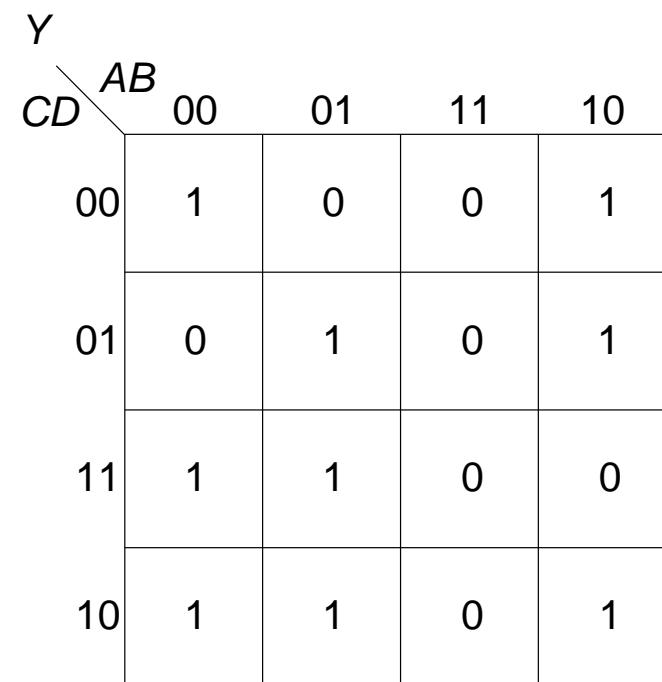
A	B	C	D	Y
0	0	0	0	1
0	0	0	1	0
0	0	1	0	1
0	0	1	1	1
0	1	0	0	0
0	1	0	1	X
0	1	1	0	1
0	1	1	1	1
1	0	0	0	1
1	0	0	1	1
1	0	1	0	X
1	0	1	1	X
1	1	0	0	X
1	1	0	1	X
1	1	1	0	X
1	1	1	1	X



$$Y = A + \bar{B}\bar{D} + C$$

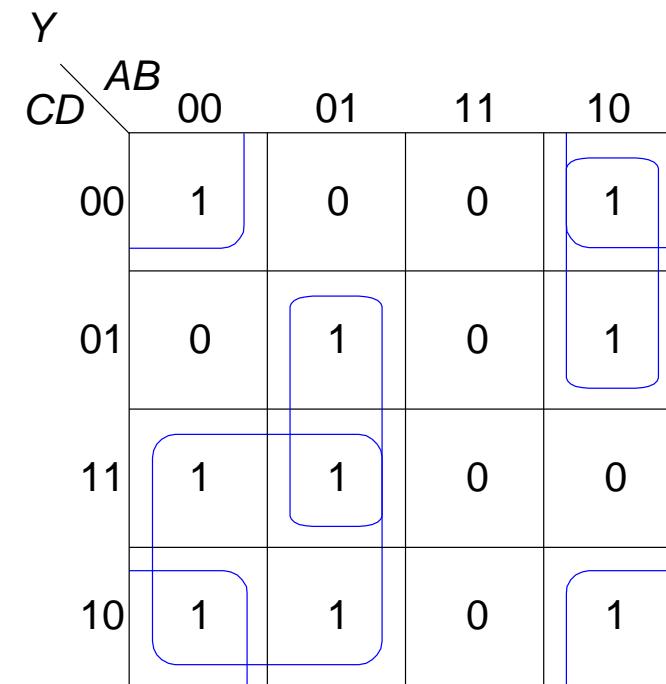
4-Input K-Map: POS & SOP Form

A	B	C	D	Y
0	0	0	0	1
0	0	0	1	0
0	0	1	0	1
0	0	1	1	1
0	1	0	0	0
0	1	0	1	1
0	1	1	0	1
0	1	1	1	1
1	0	0	0	1
1	0	0	1	1
1	0	1	0	1
1	0	1	1	0
1	1	0	0	0
1	1	0	1	0
1	1	1	0	0
1	1	1	1	0



4-Input K-Map: POS Form

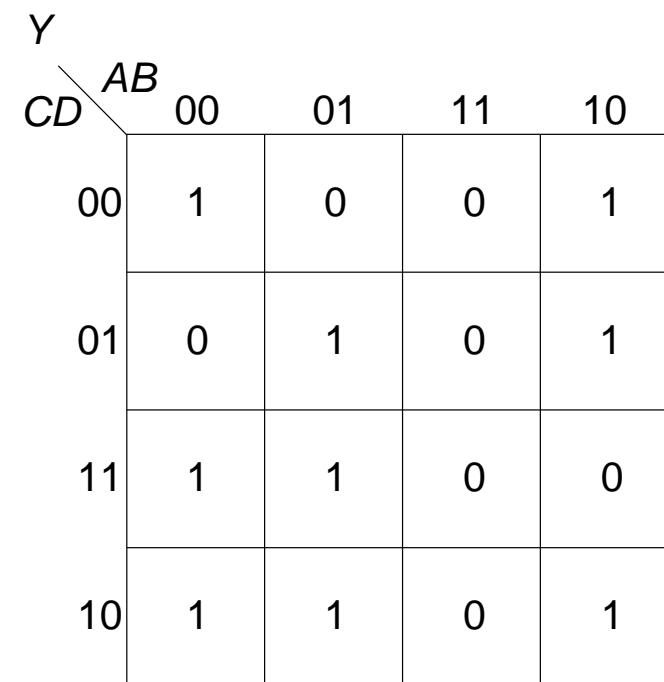
A	B	C	D	Y
0	0	0	0	1
0	0	0	1	0
0	0	1	0	1
0	0	1	1	1
0	1	0	0	0
0	1	0	1	1
0	1	1	0	1
0	1	1	1	1
1	0	0	0	1
1	0	0	1	1
1	0	1	0	1
1	0	1	1	0
1	1	0	0	0
1	1	0	1	0
1	1	1	0	0
1	1	1	1	0



$$Y = \bar{A}C + \bar{A}BD + A\bar{B}\bar{C} + \bar{B}\bar{D}$$

4-Input K-Map: SOP Form

A	B	C	D	Y
0	0	0	0	1
0	0	0	1	0
0	0	1	0	1
0	0	1	1	1
0	1	0	0	0
0	1	0	1	1
0	1	1	0	1
0	1	1	1	1
1	0	0	0	1
1	0	0	1	1
1	0	1	0	1
1	0	1	1	0
1	1	0	0	0
1	1	0	1	0
1	1	1	0	0
1	1	1	1	0



Canonical POS Expansion

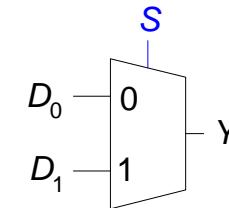
- “Add” literal/complement terms to reverse simplification (\rightarrow expand literal)
- Example
 - $Y = C$
 - $Y = C + A\bar{A}$
 - $Y = (C + A) \cdot (C + \bar{A})$
 - $Y = [(C + A) + B\bar{B}](C + \bar{A})$
 - $Y = [(C + A + B)(C + A + \bar{B})](C + \bar{A})$
 - ...

Combinational Building Blocks

- Multiplexers
- Decoders

Multiplexer (Mux)

- Selects between one of N inputs to connect to output
 - $\log_2 N$ -bit required to select input – control input S
- Example:
2:1 Mux (2 inputs to 1 output)
 - $N = 2$
 - $\log_2 2 = 1$ control bit required

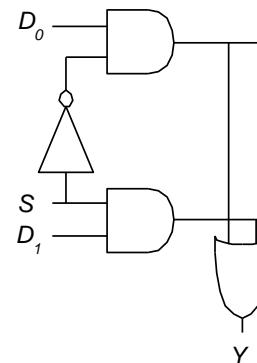
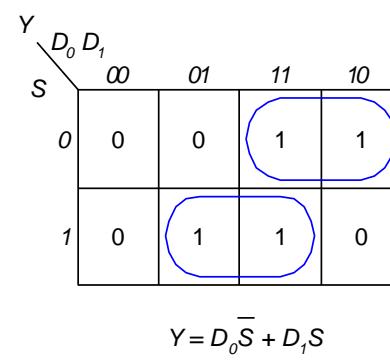


S	D_1	D_0	Y		
				0	D_0
0	0	0	0	0	D_0
0	0	1	1	1	D_1
0	1	0	0		
0	1	1	1		
1	0	0	0		
1	0	1	0		
1	1	0	1		
1	1	1	1		

Multiplexer Implementations

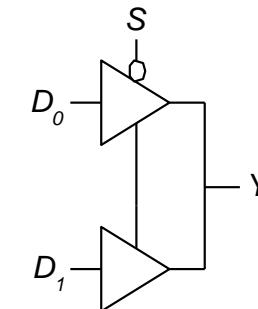
- Logic gates
 - Sum-of-products form

S	D_1	D_0	Y
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	1
1	1	1	1



- Tristates

- For an N-input mux, use N tristates
- Turn on exactly one to select the appropriate input

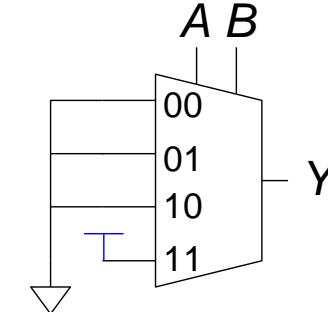


Logic using Multiplexers

- Using the mux as a lookup table
 - Zero outputs tied to GND
 - One output tied to VDD

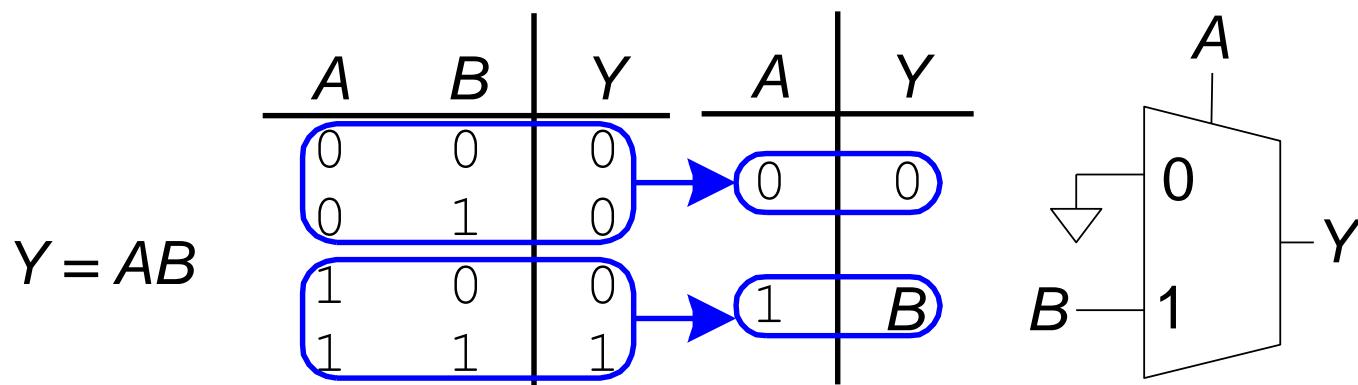
A	B	Y
0	0	0
0	1	0
1	0	0
1	1	1

$$Y = AB$$



Logic using Multiplexers

- Reducing the size of the mux

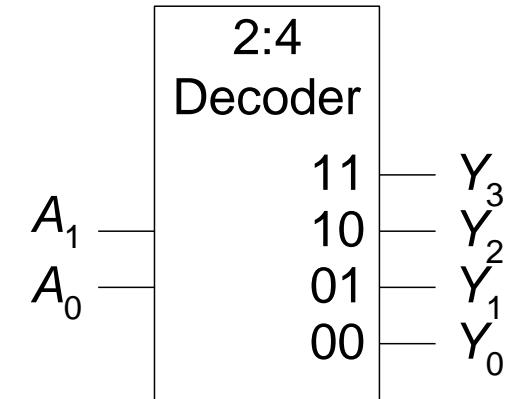


Decoders

- N inputs, 2^N outputs
- One-hot outputs: only one output HIGH at once
- Example

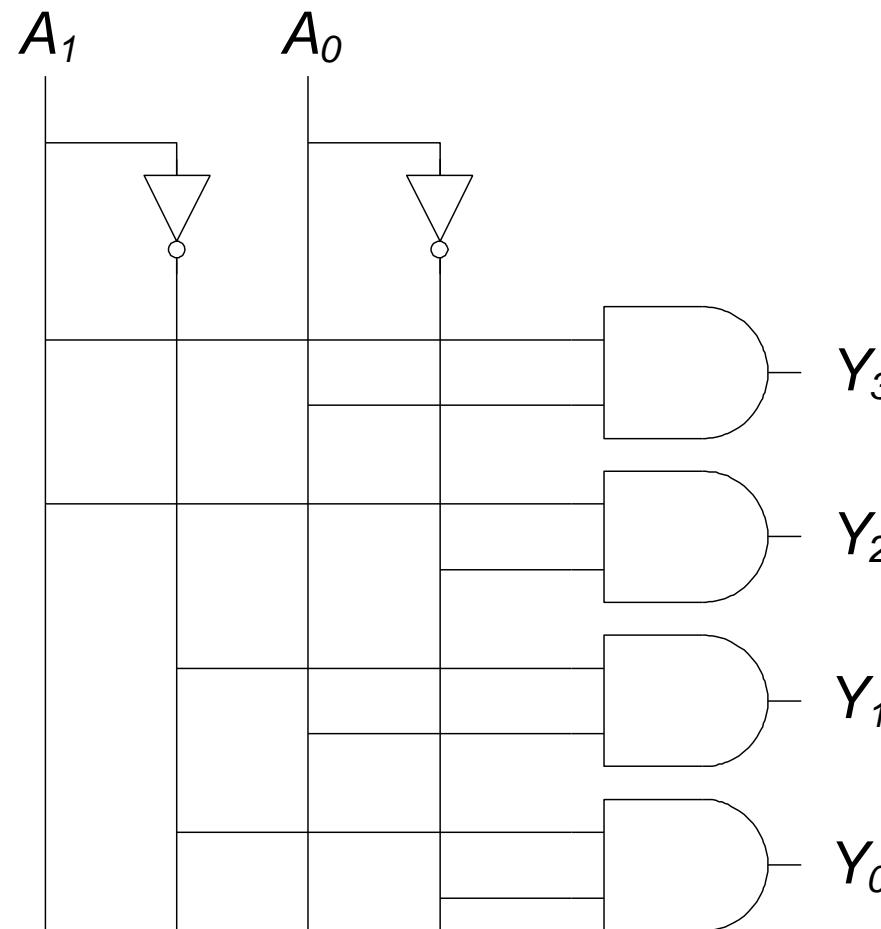
2:4 Decoder (2 inputs to 4 outputs)

- A_i decimal value selects the corresponding output



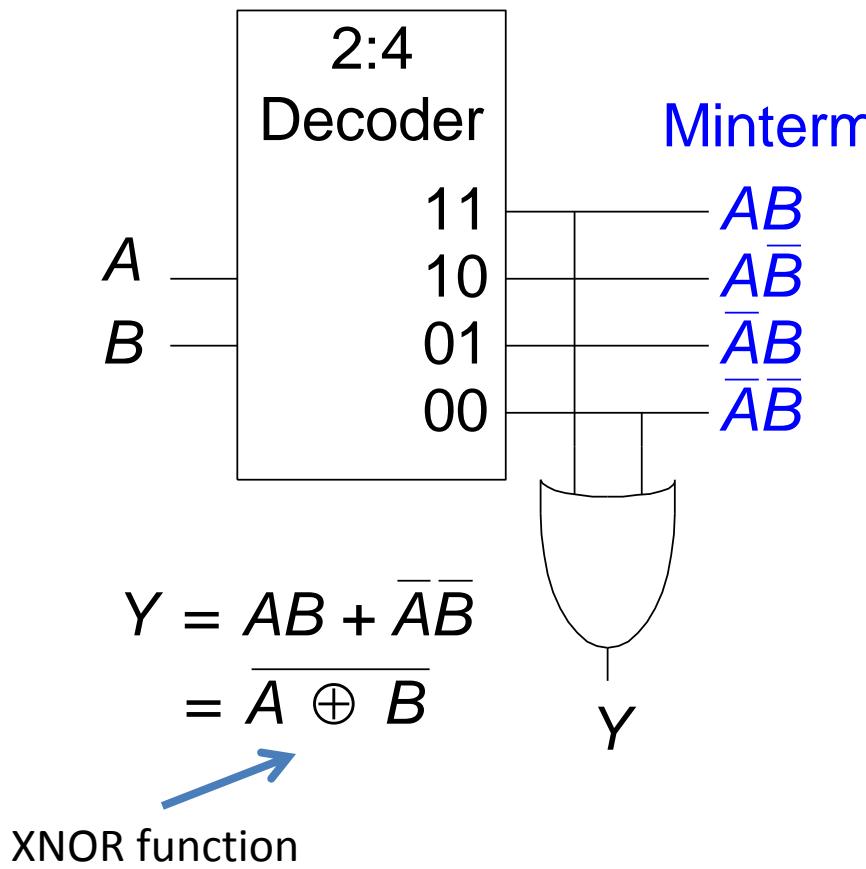
A_1	A_0	Y_3	Y_2	Y_1	Y_0
0	0	0	0	0	1
0	1	0	0	1	0
1	0	0	1	0	0
1	1	1	0	0	0

Decoder Implementation



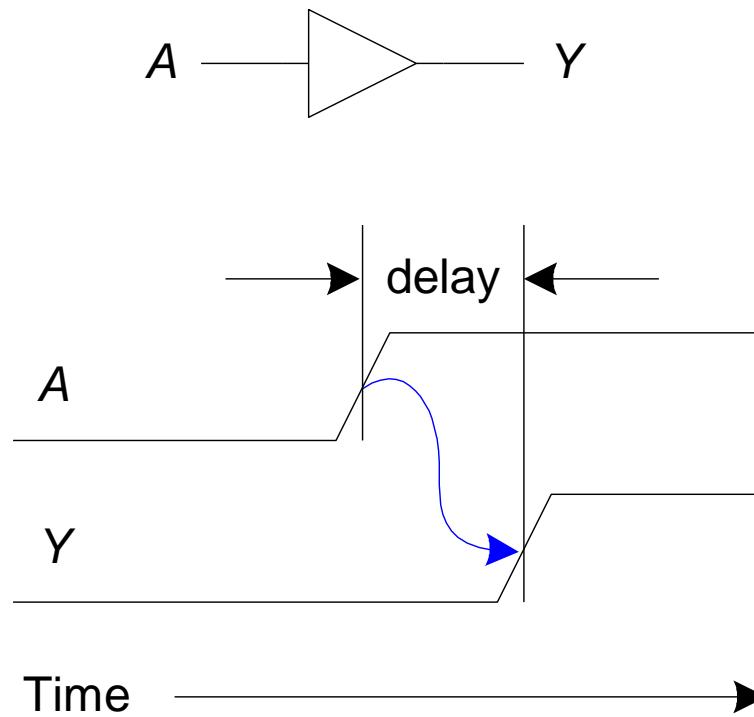
Logic Using Decoders

- OR minterms



Timing

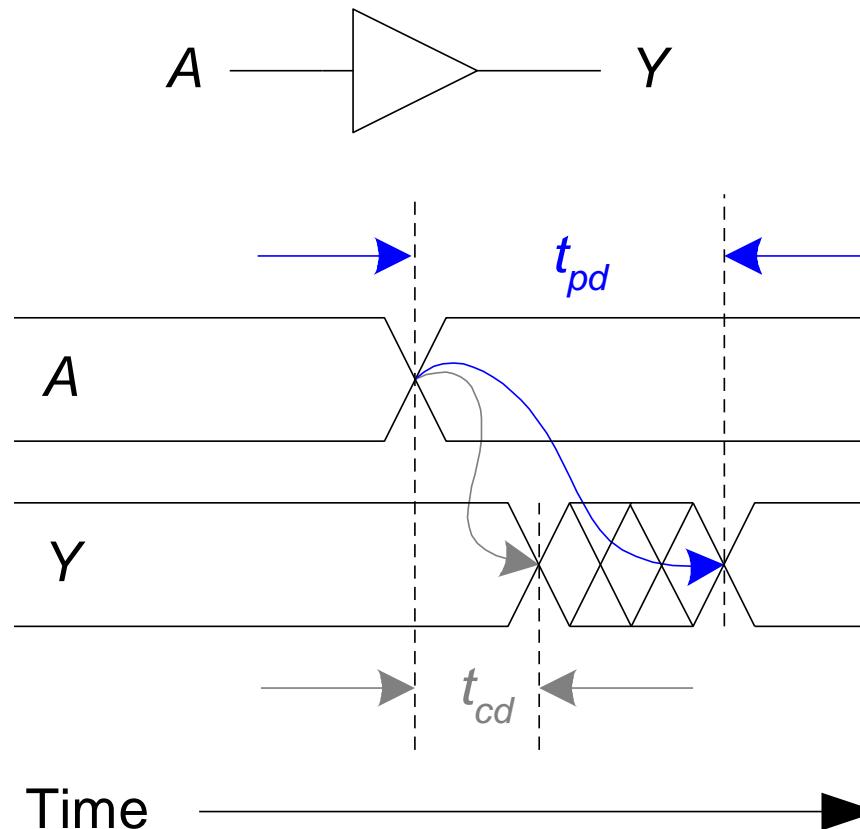
- Delay between input change and output changing



- How to build fast circuits?

Propagation & Contamination Delay

- **Propagation delay:** t_{pd} = max delay from input to final output
- **Contamination delay:** t_{cd} = min delay from input to initial output change



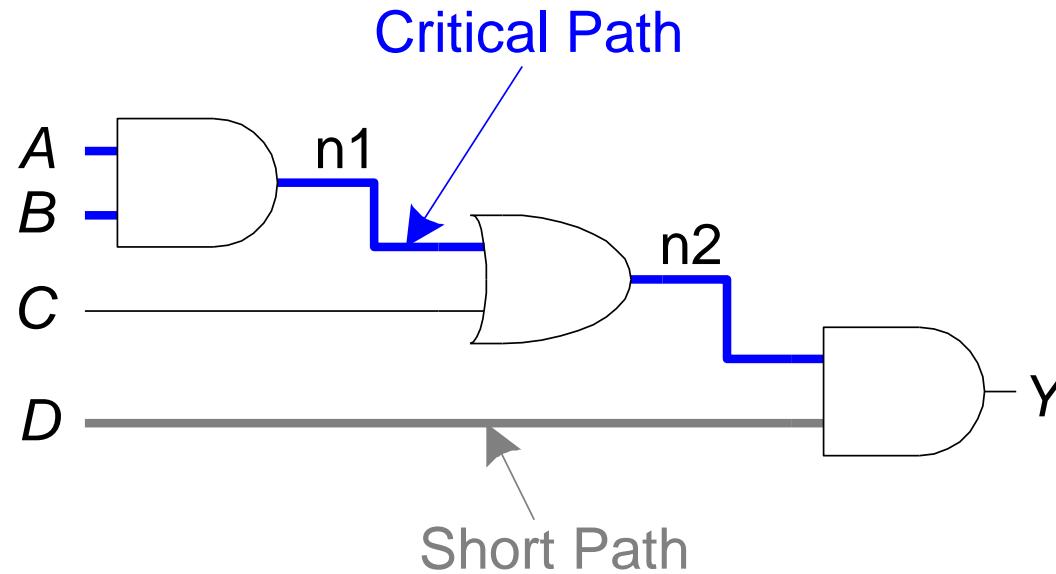
Note: Timing diagram shows a signal with a high and low and transition time as an 'X'.

Cross hatch indicates unknown/changing values

Propagation & Contamination Delay

- Delay is caused by
 - Capacitance and resistance in a circuit
 - Speed of light limitation
- Reasons why t_{pd} and t_{cd} may be different:
 - Different rising and falling delays
 - Multiple inputs and outputs, some of which are faster than others
 - Circuits slow down when hot and speed up when cold

Critical (Long) & Short Paths



Critical (Long) Path: $t_{pd} = 2t_{pd_AND} + t_{pd_OR}$

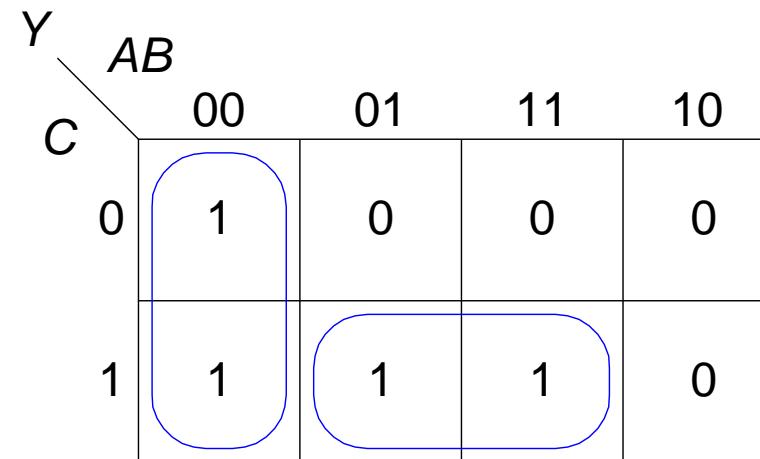
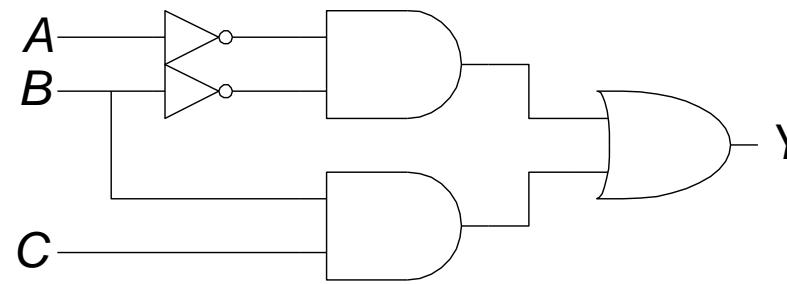
Short Path: $t_{cd} = t_{cd_AND}$

Glitches

- When a single input change causes an output to change multiple times

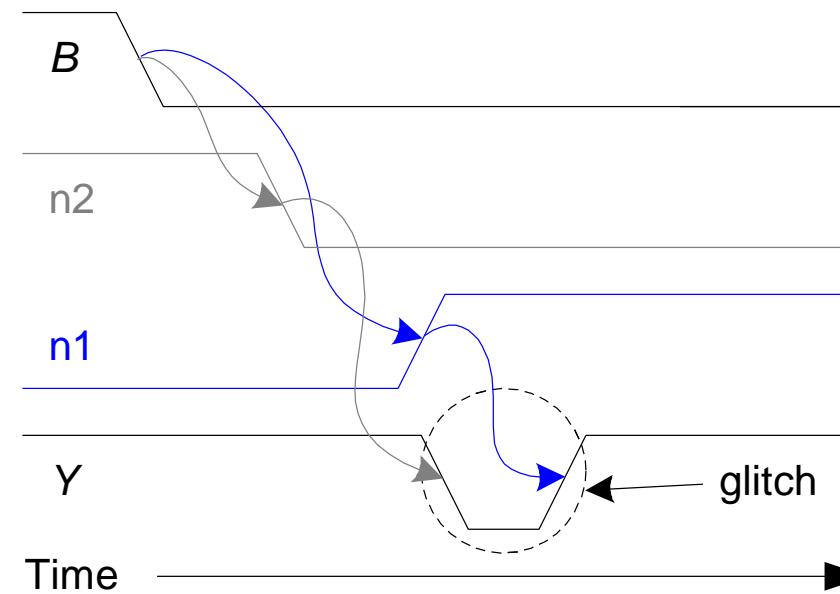
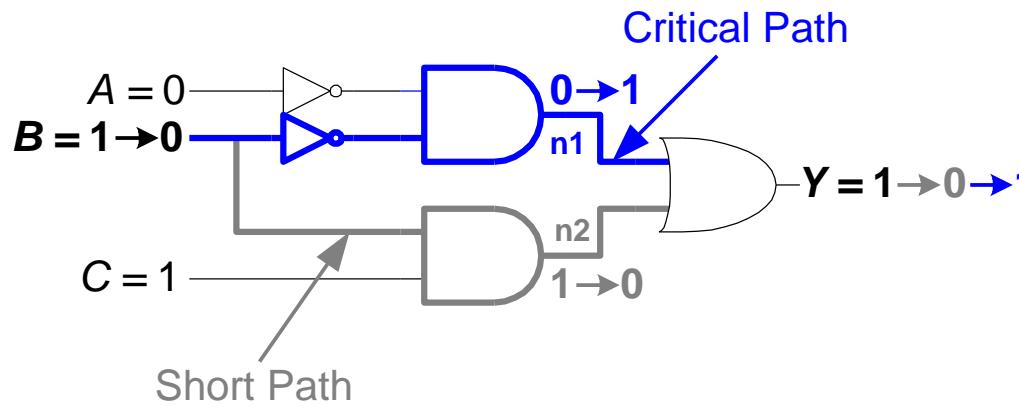
Glitch Example

- What happens when $A = 0, C = 1, B$ falls?



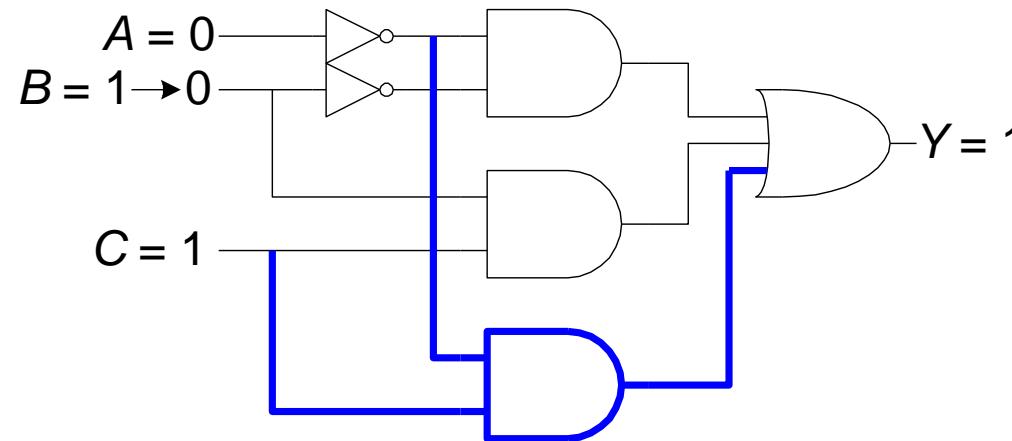
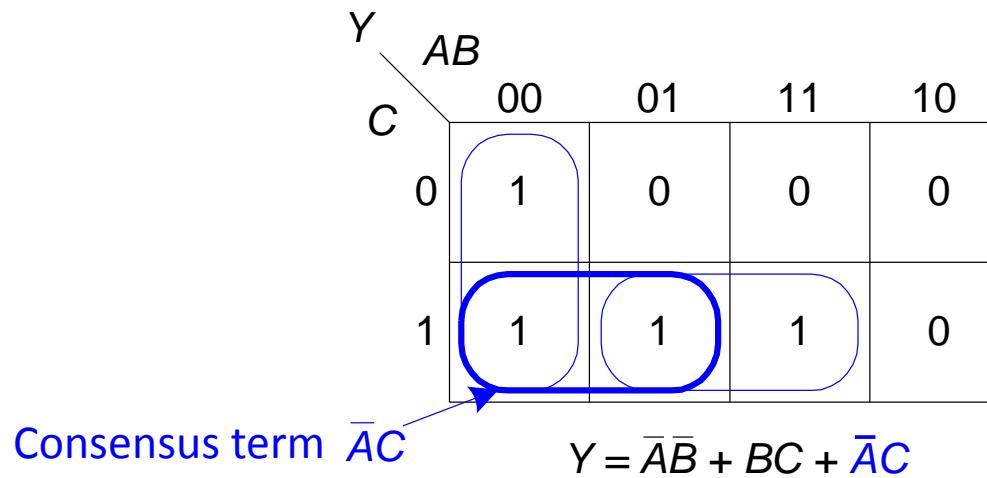
$$Y = \bar{A}\bar{B} + BC$$

Glitch Example (cont.)



Note: n_1 is slower than n_2 because of the extra inverter for B to go through

Fixing the Glitch



Why Understand Glitches?

- Glitches shouldn't cause problems because of **synchronous design** conventions (see Chapter 3)
- It's important to **recognize** a glitch: in simulations or on oscilloscope
- Can't get rid of all glitches – simultaneous transitions on multiple inputs can also cause glitches